

## 1 Introduction

*Nowhere-zero flows* were introduced by Tutte in relation to graph colorings and chromatic polynomials [6]. The concept of nowhere-zero flows (NZF for short) has connections to other graph theoretical problems such as cycle covers [3, 7] and Woodall's conjecture about dicuts and dijoins [1]. Furthermore, NZFs can also be generalized from graphs to regular matroids [5].

**Definition 1.** Let  $G = (V, E)$  be a graph and  $k$  a positive integer. We call a pair  $(D, x)$  a nowhere-zero  $k$ -flow ( $k$ -NZF) on  $G$  if

- $D = (V, A)$  is an orientation of  $G$ ,
- $x : A \rightarrow [1 - k, k - 1] \cap \mathbb{Z}$  is a circulation on  $D$ , that is,

$$\sum_{uv \in A} x(uv) - \sum_{vw \in A} x(vw) = 0 \quad \forall v \in V,$$

- $x(uv) \neq 0 \quad \forall uv \in A$ .

We can replace  $[1 - k, k - 1] \cap \mathbb{Z}$  by an abelian group  $\Gamma$  to get the definition of nowhere-zero  $\Gamma$ -flows ( $\Gamma$ -NZFs).

**Observation 1.** *The following are equivalent:*

1.  $\exists D$  orientation of  $G$  such that  $G$  admits a  $k$ -NZF with orientation  $D$ .
2.  $\forall D$  orientation,  $G$  admits a  $k$ -NZF with orientation  $D$ .
3.  $G$  admits a  $k$ -NZF  $(D, x)$ , where  $x(e) > 0 \quad \forall e \in A$ .

Sometimes it is convenient to work with finite groups instead of  $\mathbb{Z}$ , therefore we state the following theorem.

**Theorem 2** (Jaeger [2]). *Let  $G = (V, E)$  be a graph and  $\Gamma$  any finite group with  $|\Gamma| = k$ . Then  $G$  admits a  $\Gamma$ -NZF if and only if  $G$  admits a  $k$ -NZF.*

## 2 Existence of a $k$ -NZF

In planar graphs there is a clear connection between map-colorings and nowhere-zero flows captured in the following statement.

**Theorem 3.** *Let  $G = (V, E)$  be a 2-edge-connected planar graph. Then the faces of  $G$  can be colored with  $k$  colors if and only if  $G$  admits a nowhere-zero  $k$ -flow.*

By the four color theorem, we know that the faces of every loopless planar graph can be properly colored using 4 colors. Based on the analogy between map-colorings and nowhere-zero flows it makes sense to ask the question, whether we have such universal upper bound for nowhere-zero flows. Before we answer this question, we show a few special cases for  $k$ -NZFs for small values of  $k$ .

**Proposition 1.** *A graph admits a 2-NZF if and only if every vertex of  $G$  has even degree.*

**Proposition 2.** *Let  $G$  be a 3-regular graph. Then*

1.  $G$  admits a 3-NZF if and only if  $G$  is bipartite.
2.  $G$  admits a 4-NZF if and only if  $\chi'(G) = 3$ .

*Proof.* The first part is left as an exercise for the reader. For the second part, by Theorem 2, it suffices to prove the statement for  $\mathbb{Z}_2 \times \mathbb{Z}_2$ -NZF instead of 4-NZF. Finishing the proof is an easy exercise left for the reader.  $\square$

To answer our question about the universal upper bound for a  $k$ -NZF, we first state Tutte's conjecture.

**Conjecture 1** (Tutte, 1954 [6]). *Every bridge-less graph admits a 5-NZF.*

The condition that  $G$  is bridge-less is clearly necessary since a graph that has a bridge cannot have a  $k$ -NZF for any  $k$ . The next observation shows that 5 is the lowest upper bound we can hope for.

**Observation 2.** *The Petersen graph does not admit a 4-NZF.*

*Proof.* The edges of the Petersen graph cannot be colored with 3 colors, thus by Proposition 2 it cannot admit a 4-NZF.  $\square$

**Conjecture 2.** *If a bridge-less graph does not contain a subdivided Petersen graph, then it admits a 4-NZF.*

The first universal upper bound was proven by Jaeger in 1979.

**Theorem 4** (Jaeger [2]). *Every bridge-less graph admits an 8-NZF.*

In the next section we show Seymour's improvement on this result to 6, which is the best known upper bound.

### 3 The main result

**Theorem 5** (Seymour [4]). *Every bridge-less graph admits a 6-NZF.*

First we need the following observations:

**Proposition 3.** *For  $k > 2$ , if  $G = (V, E)$  is a bridge-less graph such that  $G$  does not admit a  $k$ -NZF and  $|V| + |E|$  is minimum, then  $G$  is simple, 3-regular and 3-connected.*

**Proposition 4.** *If  $k > 0$  and  $\Phi$  is an integer-valued circulation in  $G = (V, E)$ , then there is a  $k$ -flow  $\Phi'$  such that for every  $e \in E$ ,*

$$\Phi(e) \equiv \Phi'(e) \pmod{k}.$$

To prove the main theorem, we need to introduce a closure operator on subsets of the edges:

**Definition 6.** For  $G = (V, E)$  and an integer  $k > 0$ , for a subset  $X \subseteq E$  we define  $\langle X \rangle_k$  to be the smallest set  $Y \subseteq E$  with the following properties:

- $X \subseteq Y$ ,
- there is no cycle  $C$  such that  $0 < |C - Y| \leq k$ .

It can easily be seen that  $\langle X \rangle_k$  is well-defined because if  $Y_1$  and  $Y_2$  satisfy the above properties, then so does  $Y_1 \cap Y_2$ . An important property of  $\langle X \rangle_k$  is that it is a closure operator, that is,

$$X \subseteq \langle X \rangle_k, \quad \langle \langle X \rangle_k \rangle_k = \langle X \rangle_k, \quad X \subseteq Y \Rightarrow \langle X \rangle_k \subseteq \langle Y \rangle_k.$$

**Lemma 7.** *Let  $G = (V, E)$ ,  $k > 0$ , let  $X \subseteq E$  be such that  $\langle X \rangle_{k-1} = E$ . Then  $G$  admits a  $k$ -flow  $\Phi$  with  $E - X \subseteq \text{supp}(\Phi)$ .*

*Proof.* We proceed by induction on  $|E - X|$ . If this is zero the result is trivial, so we assume it is not the case. Then  $\langle X \rangle_{k-1} \neq X$  and so there is a circuit  $C$  with  $0 < |C - X| \leq k - 1$ . By the properties of  $\langle \cdot \rangle_k$ ,  $\langle XUC \rangle_{k-1} = E$ , and so by induction there is a  $k$ -flow  $\Phi_1$  with  $E - (XUC) \subseteq \text{supp}(\Phi_1)$ . Take a circulation  $\Phi_2$  so that  $\text{supp}(\Phi_2) = C$  and  $\Phi_2(e) = 0 \forall e \in E - C, \Phi_2(e) = \pm 1 \forall e \in C$ . Choose an integer  $\ell$  such that  $0 \leq \ell \leq k - 1$  and  $\Phi_1(e) + \ell\Phi_2(e) \neq 0$  ( $k$ ) for each  $e \in C - X$  (this is possible since  $|C - X| \leq k - 1$ ). Put  $\Phi = \Phi_1 + \ell\Phi_2$ . Then for  $e \in E - (XUC)$ ,  $\Phi(e) = \Phi_1(e)$  and for  $e \in C - X$ ,  $\Phi(e) = \Phi_1(e) + \ell\Phi_2(e) \neq 0$  ( $k$ ) by choice of  $\ell$ . Thus for each  $e \in E - X$ ,  $\Phi(e) \neq 0$  ( $k$ ). The result follows from Proposition 4.  $\square$

**Lemma 8.** *If  $G = (V, E)$  is a 3-connected graph and  $|V| \geq 3$ , then there are vertex-disjoint circuits  $C_1, \dots, C_r$  such that  $\langle C_1 \cup \dots \cup C_r \rangle_2 = E$ .*

A key idea of proving the lemma is to prove the following:

**Lemma 9.** *Let  $G = (V, E)$  be a simple graph with  $d(v) \geq 2$  for every vertex  $v \in V$ . Then there is a subgraph  $B$  of  $G$  with at least 3 vertices, such that  $B$  is 2-connected and at most one vertex of  $B$  is adjacent in  $G$  to vertices of  $G$  not in  $B$ .*

We do not show the proof of either lemma here, they can be found in [4].

Now we finally prove the main theorem using this lemma.

*Proof of Theorem 5.* By Proposition 3, it is enough to prove the theorem for 3-regular, simple, 3-connected graphs. In this case, by Lemma 8, there exists a subset  $X \subseteq E$  such that  $X = C_1 \cup \dots \cup C_r$  for some vertex-disjoint circuits  $C_1 \dots C_r$  and  $\langle X \rangle_2 = E$ . By Lemma 7, there is a 3-flow  $\Phi_1$  such that  $E - X \subseteq \text{supp}(\Phi_1)$ . Let  $\Phi_2$  be a 2-flow on  $X$ . Then  $\Phi = \Phi_1 + 3\Phi_2$  is clearly a flow. For an edge  $e \in E - X$ ,  $\Phi(e) = \Phi_1(e) \neq 0$  and  $|\Phi(e)| \leq 2$ . For an edge  $e \in X$ ,  $\Phi(e) = \Phi_1(e) \pm 3 \neq 0$  and  $|\Phi(e)| \leq 5$ , thus  $\Phi$  is a nowhere-zero 6-flow.  $\square$

## References

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