

Optimization

Fall semester 2022/23

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Lydia Mendoza

Eötvös Loránd University
Institute of Mathematics
Department of Operations Research



Set 1

Question 1

Bob would like to write down the system

$$3x + 2y + 4z = 8, \quad (1)$$

$$-3y \leq 3, \quad (2)$$

$$x - 3z \geq 10, \quad (3)$$

$$\min x - y, \quad (4)$$

but his keyboard is missing the symbols $=$ and \geq , and the letter i is not working.

Reformulate the problem only using \leq and maximization. (1pt)

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Recall. We can write $3x + 2y + 4z = 8$ using the two inequalities $3x + 2y + 4z \leq 8$ and $3x + 2y + 4z \geq 8$.

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Also, $Ax' = A \frac{x}{\min_i x_i} \leq 0$ as $Ax \leq 0$ and $\frac{1}{\min_i x_i} \geq 0$

Question 3

Consider the problem

$$x_2 \leq 4, \quad (5)$$

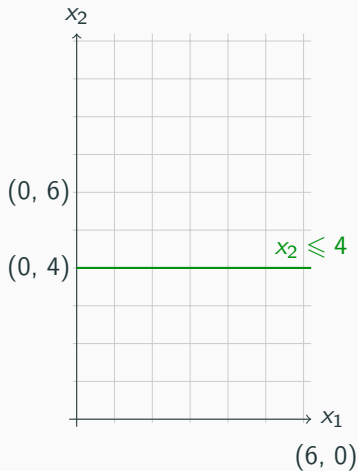
$$x_1 + x_2 \leq 6, \quad (6)$$

$$2x_1 + x_2 \leq 10 \quad (7)$$

$$x_1, x_2 \geq 0 \quad (8)$$

Represent these constraints on the plane.

Find a point that maximizes $x_1 + 2x_2$. (2pts)



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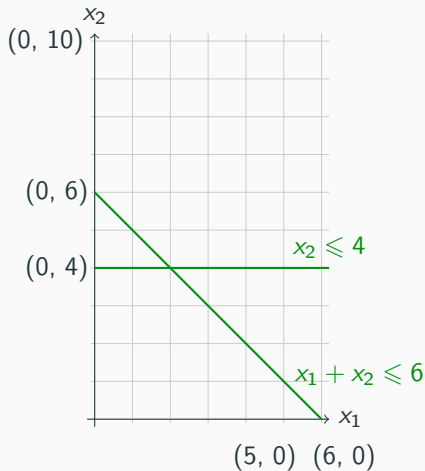
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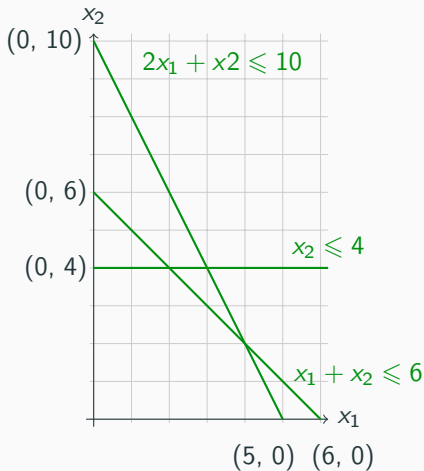
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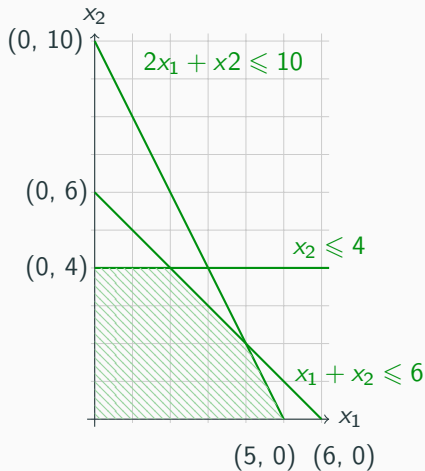
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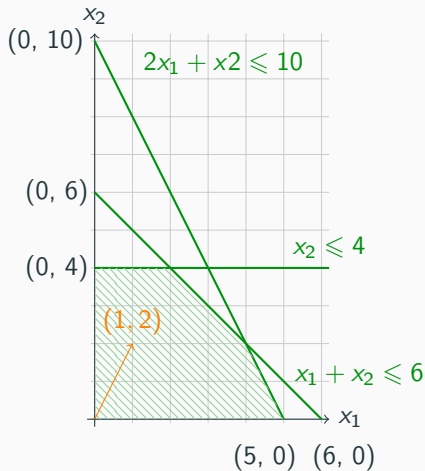
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Optimal solution: $x_1 = 2, x_2 = 4$.
Optimal value: $x_1 + 2x_2 = 2 + 2 * 4 = 10$.

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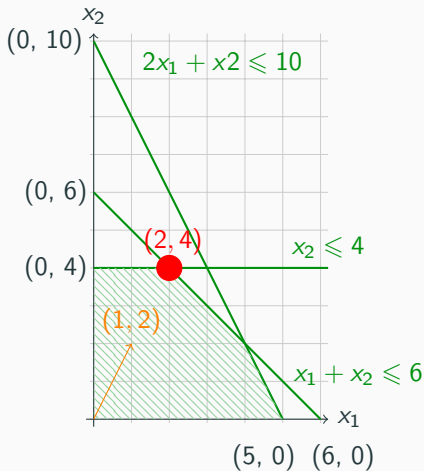
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Question 3: Strong Duality Thm

Consider the problem

$$x_2 \leq 4, \quad (9)$$

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Optimal primal solution: $x_1 = 2, x_2 = 4$.

Optimal primal value: $x_1 + 2x_2 = 2 + 2 * 4 = 10$.

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The dual of this problem is:

$$\min 4y_1 + 6y_2 + 10y_3$$

$$\text{s.t.} \quad y_2 + 2y_3 \geq 1$$

$$y_1 + y_2 + y_3 \geq 2$$

$$y_1, y_2, y_3 \geq 0$$

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Optimal dual value = 10.

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max = min as the Strong Duality theorem states.

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Verify the 'only if' direction in the general form of Farkas' lemma. (1pt)

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We want to show that

If there is no $y = (y_0, y_1)$ s.t.

$$y_0P + y_1Q = 0 \quad (13)$$

$$y_0A + y_1B \geq 0 \quad (14)$$

$$y_1 \geq 0 \quad (15)$$

$$y_0b_0 + y_1b_1 < 0 \quad (16)$$

then

there exists $x = (x_0, x_1)$ s.t.

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$x = (x_0, x_1)$ and $y = (y_0, y_1)$ exists.

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$$\begin{aligned} 0 &> y_0b_0 + y_1b_1 \\ &= y_0(Px_0 + Ax_1) + y_1b_1 \end{aligned}$$

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$x = (x_0, x_1)$ and $y = (y_0, y_1)$ exists.

Suppose by the contrary that both

$x = (x_0, x_1)$ and $y = (y_0, y_1)$ exists.

Then

$$0 > y_0b_0 + y_1b_1$$

$$= y_0(Px_0 + Ax_1) + y_1b_1$$

$$\geq y_0(Px_0 + Ax_1) + y_1(Qx_0 + Bx_1)$$

$$= (y_0P + y_1Q)x_0 + (y_0A + y_1B)x_1$$

$$= 0 + (y_0A + y_1B)x_1$$

$$\geq 0,$$

which is a contradiction.

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Dual:

$$\min y_0b_0 + y_1b_1 \quad (24)$$

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Question 6

Let $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$, and $c_1, \dots, c_k \in \mathbb{R}^n$. Formulate the following problem as an LP: $Ax = b, x \geq 0, \min f(x)$ where $f(x) := \max\{c_1x, \dots, c_kx\}$. (1pt)

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$$\min z$$

Reduce the following systems of inequalities to each other (in the sense that if we can solve one of them, then we can solve any of them):

(I): $Ax = b$

$x \geq 0$

(II): $Bx \leq b$

$x \geq 0$

(III): $Qx \leq b$

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Write up Farkas' lemma for all of them. (3pts)

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Let us define $B' := [Q \quad -Q]$, $b' = b$, and $x' = (x^+, x^-) \geq 0$.

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We know how to solve (III), and we would like to solve (IV).

Let us define $Q' := \begin{bmatrix} P & 0 \\ -P & 0 \\ 0 & Q \end{bmatrix}$, $b = (b_0, -b_0, b_1)$, and $x' = (x_0, x_1)$.

$$\text{(I): } Ax = b \\ x \geq 0$$

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We know how to solve (IV), and we would like to solve (I).

Let us define $P' := A$, $Q' = -Id$, $b'_0 = b$, $b'_1 = 0$, and $x'_0 = x$, $x'_1 = x$.

$$\text{(I): } Ax = b \\ x \geq 0$$

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We know how to solve (III), and we would like to solve (IV).

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We know how to solve (IV), and we would like to solve (I).

Let us define $P' := A$, $Q' = -Id$, $b'_0 = b$, $b'_1 = 0$, and $x'_0 = x$, $x'_1 = x$.

Then $P'x'_0 = b'_0$, $Q'x'_1 \leq b'_1$, is equivalent to (I).

We write the Farkas' Lemma for each problem:

$$\begin{array}{ll} \text{(I). There exists } x \text{ s.t. } Ax = b & \iff \text{there is no } y \text{ s.t. } yA \geq 0 \\ & x \geq 0 & yb < 0 \end{array}$$

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$$\text{(IV). There exists } x \text{ s.t. } \begin{array}{l} Px_0 = b_0 \\ Qx_1 \leq b_1 \end{array} \iff \begin{array}{l} \text{there is no } y \text{ s.t. } y_0P + y_1Q = 0 \\ y_1 \geq 0 \\ y_0b_0 + y_1b_1 < 0 \end{array}$$